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Online

# Guaranteeing Timed Opacity using Parametric Timed Model Checking

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## Context: timing attacks

- ▶ Principle: deduce **private information** from timing data (**execution time**)

Issues:

- ▶ May depend on the **implementation** (or, even worse, be **introduced by the compiler**)
- ▶ A relatively trivial solution: make the program last always its maximum execution time  
Drawback: **loss of efficiency**

↪ Non-trivial problem

## A simple example of timing attack

```
1 # input pwd      : Real password
2 # input attempt: Tentative password
3 for i = 0 to min(len(pwd), len(attempt)) - 1 do
4     if pwd[i] /= attempt[i] then
5         return false
6 done
7 return true
8
```

## A simple example of timing attack

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1 # input pwd      : Real password
2 # input attempt: Tentative password
3 for i = 0 to min(len(pwd), len(attempt)) - 1 do
4     if pwd[i] != attempt[i] then
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pwd	c	h	i	c	k	e	n
-----	---	---	---	---	---	---	---

attempt	c	h	e	e	s	e
---------	---	---	---	---	---	---

Execution time:

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Execution time:  $\epsilon$

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Execution time:  $\epsilon + \epsilon + \epsilon$

- **Problem:** The execution time is proportional to the number of consecutive correct characters from the beginning of attempt



## Informal problems

Question: can we exhibit **secure execution times**?

### Time-opacity computation

Exhibit **execution times** for which it is not possible to infer information on the internal behavior

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### Time-opacity computation

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Further question: can we also tune internal timing constants to make the system resisting to timing attacks?

### Time-opacity synthesis

Exhibit **execution times and internal timing constants** for which it is not possible to infer information on the internal behavior

# Outline

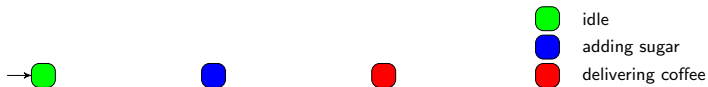
Formalism and Computation results

Toward parameter synthesis

Perspectives

# Timed automaton (TA)

- ▶ Finite state automaton (sets of *locations*)



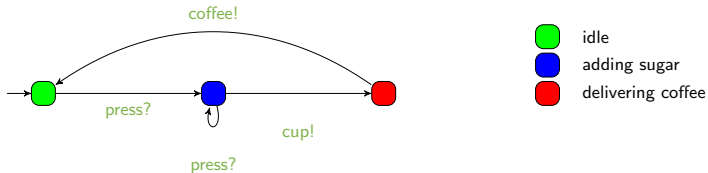
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[AD94] Rajeev Alur and David L. Dill. "A theory of timed automata". In: *Theoretical Computer Science*

126.2 (Apr. 1994), pp. 183–235. ISSN: 0304-3975. DOI: 10.1016/0304-3975(94)90010-8

# Timed automaton (TA)

- ▶ Finite state automaton (sets of **locations** and **actions**)

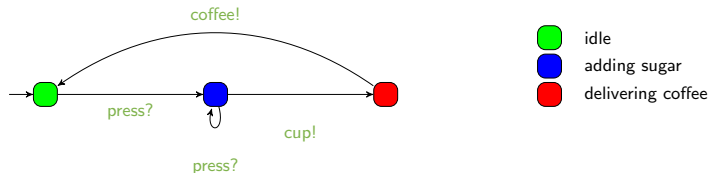


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## Timed automaton (TA)

- ▶ Finite state automaton (sets of **locations** and **actions**) augmented with a set  $X$  of **clocks** [AD94]
  - ▶ Real-valued variables evolving linearly **at the same rate**

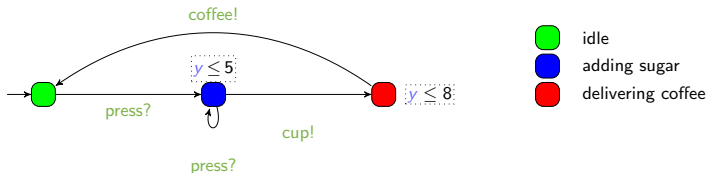


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  - ▶ Can be compared to integer constants in invariants
- ▶ Features
  - ▶ Location **invariant**: property to be verified to stay at a location

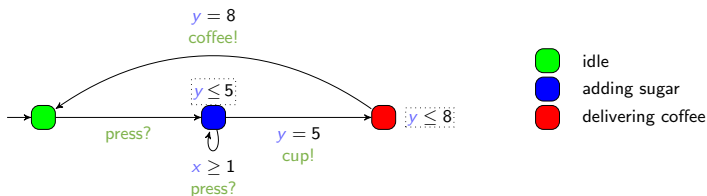


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  - ▶ Transition **guard**: property to be verified to enable a transition



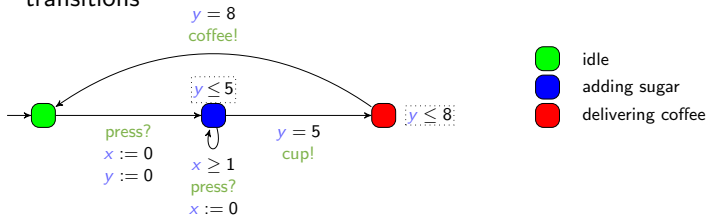
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- ▶ Features
  - ▶ Location **invariant**: property to be verified to stay at a location
  - ▶ Transition **guard**: property to be verified to enable a transition
  - ▶ Clock **reset**: some of the clocks can be **set to 0** along transitions



[AD94] Rajeev Alur and David L. Dill. "A theory of timed automata". In: *Theoretical Computer Science*

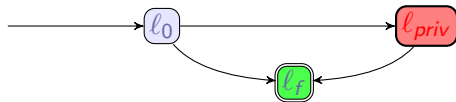
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## Formalization

Hypotheses:

[AS19]

- ▶ A start location  $l_0$  and an end location  $l_f$
- ▶ A special private location  $l_{priv}$



### Definition (timed opacity)

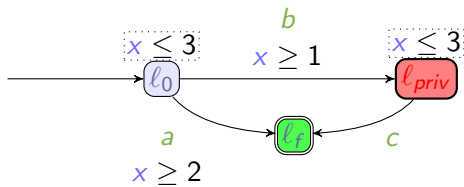
The system is **opaque** w.r.t.  $l_{priv}$  on the way to  $l_f$  for a **duration  $d$**  if there exist two runs to  $l_f$  of duration  $d$

1. one passing by  $l_{priv}$
2. one *not* passing by  $l_{priv}$

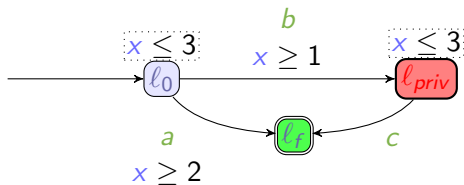
[AS19] Étienne André and Jun Sun. "Parametric Timed Model Checking for Guaranteeing Timed Opacity".

In: ATVA (Oct. 28–31, 2019). Ed. by Yu-Fang Chen, Chih-Hong Cheng, and Javier Esparza. Vol. 11781. Lecture Notes in Computer Science. Taipei, Taiwan: Springer, 2019, pp. 115–130. DOI: 10.1007/978-3-030-31784-3\_7

## Example

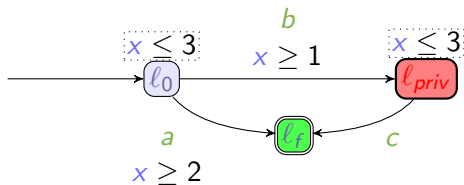


## Example



- There exist two runs of duration  $d = 2$ :

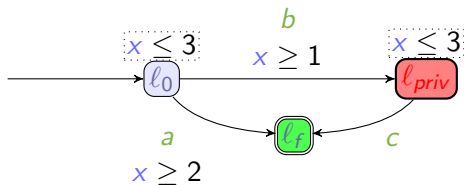
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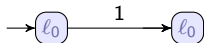
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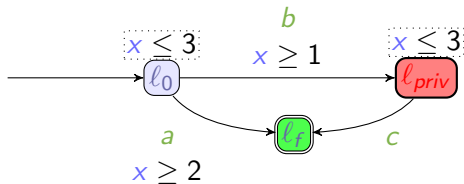
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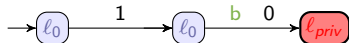
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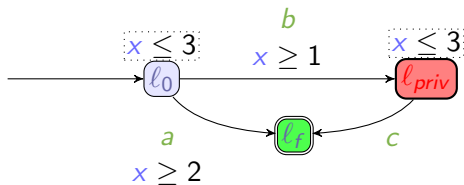
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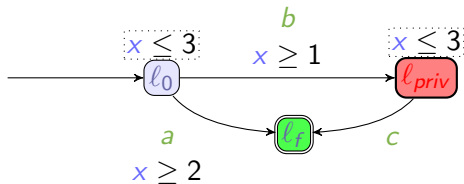


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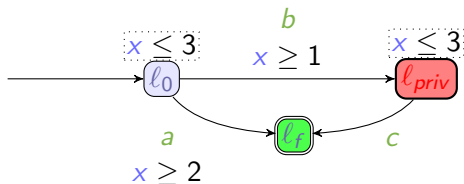
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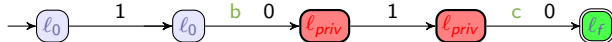
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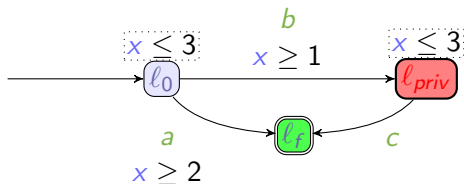
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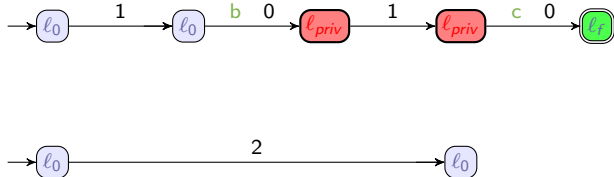
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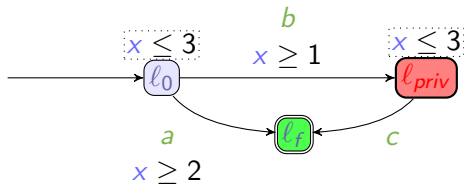
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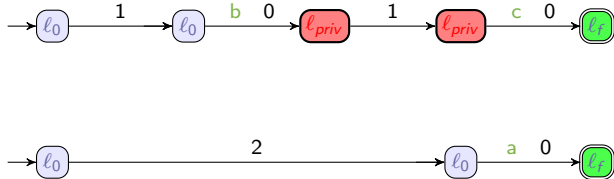
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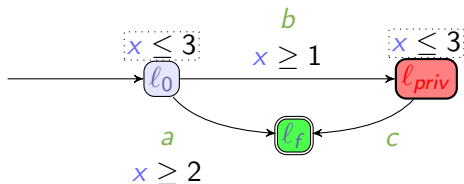
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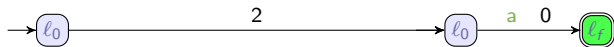
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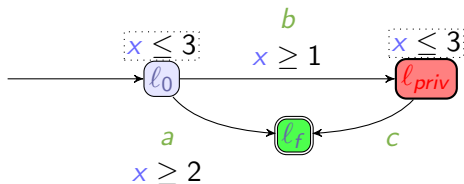


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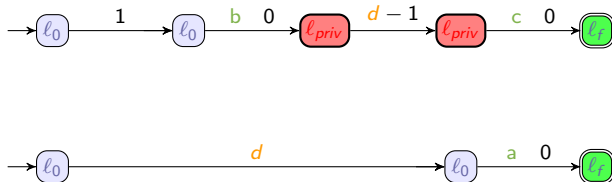


We say that the system is **opaque** w.r.t.  $l_{priv}$  on the way to  $l_f$  for a duration  $d = 2$

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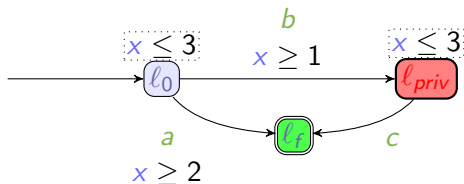


- There exist two runs of duration  $d$  for all durations  $d \in [2, 3]$ :



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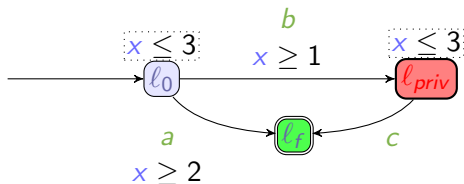
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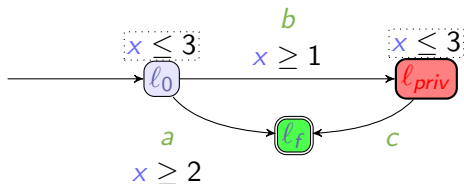
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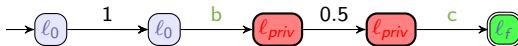


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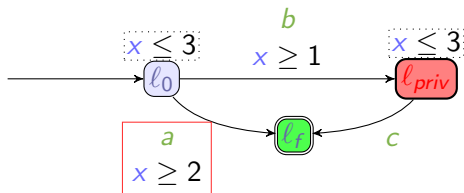
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There exists a run of duration 1.5 passing by  $l_{priv}$



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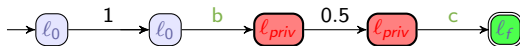


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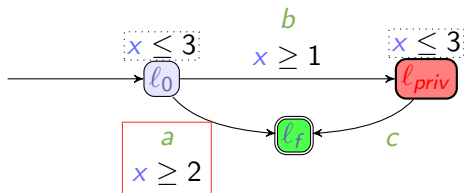
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It is not possible to reach  $l_f$  with a path of duration 1.5 not passing by  $l_{priv}$

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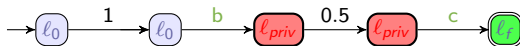


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There exists a run of duration 1.5 passing by  $l_{priv}$



It is not possible to reach  $l_f$  with a path of duration 1.5 not passing by  $l_{priv}$

We say that the system is **not fully opaque** w.r.t.  $l_{priv}$  on the way to  $l_f$

## Problem 1: timed-opacity computation

**Theorem** The durations  $d$  such that the system is opaque can be effectively computed and defined

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[And+] Étienne André, Didier Lime, Dylan Marinho, and Jun Sun. “Guaranteeing Timed Opacity using Parametric Timed Model Checking”. *submitted*

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**Theorem** The durations  $d$  such that the system is opaque can be effectively computed and defined

**Corollary** Asking if a TA is opaque for all its execution times is decidable

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## Problem 1: timed-opacity computation

**Theorem** The durations  $d$  such that the system is opaque can be effectively computed and defined

**Corollary** Asking if a TA is opaque for all its execution times is decidable

Proof: based on the region graph and RA-arithmetic (see [And+])

Exact complexity: unproved (EXPSpace upper bound proved, but exponential hardness seems likely)

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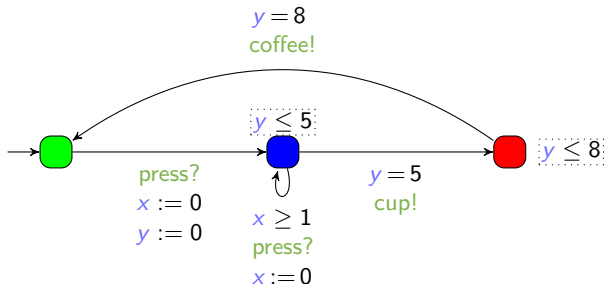
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# Parametric Timed Automaton (PTA)

- Timed automaton (sets of **locations**, **actions** and **clocks**)



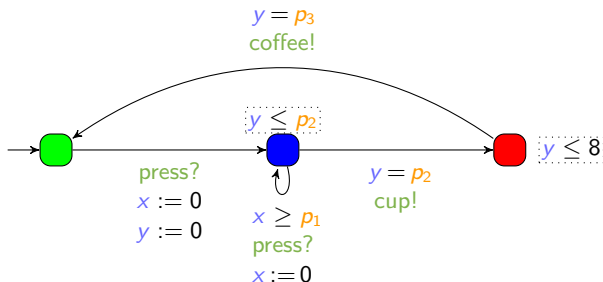

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[AHV93] Rajeev Alur, Thomas A. Henzinger, and Moshe Y. Vardi. "Parametric real-time reasoning". In: *STOC*. ed. by S. Rao Kosaraju, David S. Johnson, and Alok Aggarwal. San Diego, California, United States: ACM, 1993, pp. 592–601. ISBN: 0-89791-591-7. DOI: 10.1145/167088.167242



# Parametric Timed Automaton (PTA)

- ▶ Timed automaton (sets of **locations**, **actions** and **clocks**) augmented with a set  $P$  of **parameters** [AHV93]
  - ▶ **Unknown constants** compared to a **clock** in guards and invariants



[AHV93] Rajeev Alur, Thomas A. Henzinger, and Moshe Y. Vardi. "Parametric real-time reasoning". In: *STOC*. ed. by S. Rao Kosaraju, David S. Johnson, and Alok Aggarwal. San Diego, California, United States: ACM, 1993, pp. 592–601. ISBN: 0-89791-591-7. DOI: 10.1145/167088.167242

## Overview of our theoretical results

- ▶ General case: The mere existence of a parameter valuation for which there exists a duration for which timed-opacity is achieved **is undecidable**

[And+]

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- ▶ Study of a subclass known for being “at the frontier” of decidability (L/U-PTA)
  - ▶ The existence of valuations for timed opacity w.r.t. some execution times is decidable
  - ▶ The existence of valuations for full timed opacity is undecidable
  - ▶ The synthesis is uncomputable in practice

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[And+]

We adopt a “best-effort” approach for the general case of PTAs

- ▶ Approach not guaranteed to terminate in theory

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Formalism and Computation results

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- ▶ Some restricted problems remain open  
e. g., PTA with one clock
- ▶ Study more restrictive sub-classes, with the hope to exhibit a decidable one

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### On the practical side

- ▶ Have an automatic translation of programs to PTAs  
→ Some experiments were done, but on Java programs manually translated to PTAs
- ▶ Repairing a non-opaque system  
→ Preliminary ideas in [And+]<sup>a</sup>, but not fixed

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<sup>a</sup>[And+] Étienne André, Didier Lime, Dylan Marinho, and Jun Sun. “Guaranteeing Timed Opacity using Parametric Timed Model Checking”. [submitted](#)

## References I



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Étienne André, Didier Lime, Dylan Marinho, and Jun Sun. “Guaranteeing Timed Opacity using Parametric Timed Model Checking”. *submitted*.



## References II



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